

Optimal Standby Virtual Routers Selection for Node Failures in a Virtual Network Environment

Introduction

Network virtualization allows flexibility to configure virtual networks in a dynamic manner. In such a setting, to provide resilient services to virtual networks, we consider the situation where the substrate network provider wants to have standby virtual routers ready to serve the virtual networks in the event of a failure. Such a failure can affect one or more virtual routers in multiple virtual networks. The goal of our work is to make the optimal selection of standby virtual routers so that virtual networks can be dynamically reconfigured back to their original topologies after a failure. We present an optimization formulation and a preliminary implementation on GENI testbed by applying the idea behind the model.

Dynamic Reconfiguration Scheme For Virtual Networks

- Substrate network providers are capable to provide standby virtual routers (S-VRs) when leasing virtual networks to a service provider.
- Requirements to the reconfiguration
 - Conservation
 - Topological structure
 - Bandwidth requirement
 - Fast restoration
 - Focus on the core node failures
- Software-defined management for reconfiguration
 - Autonomic Management
 - Central Controller

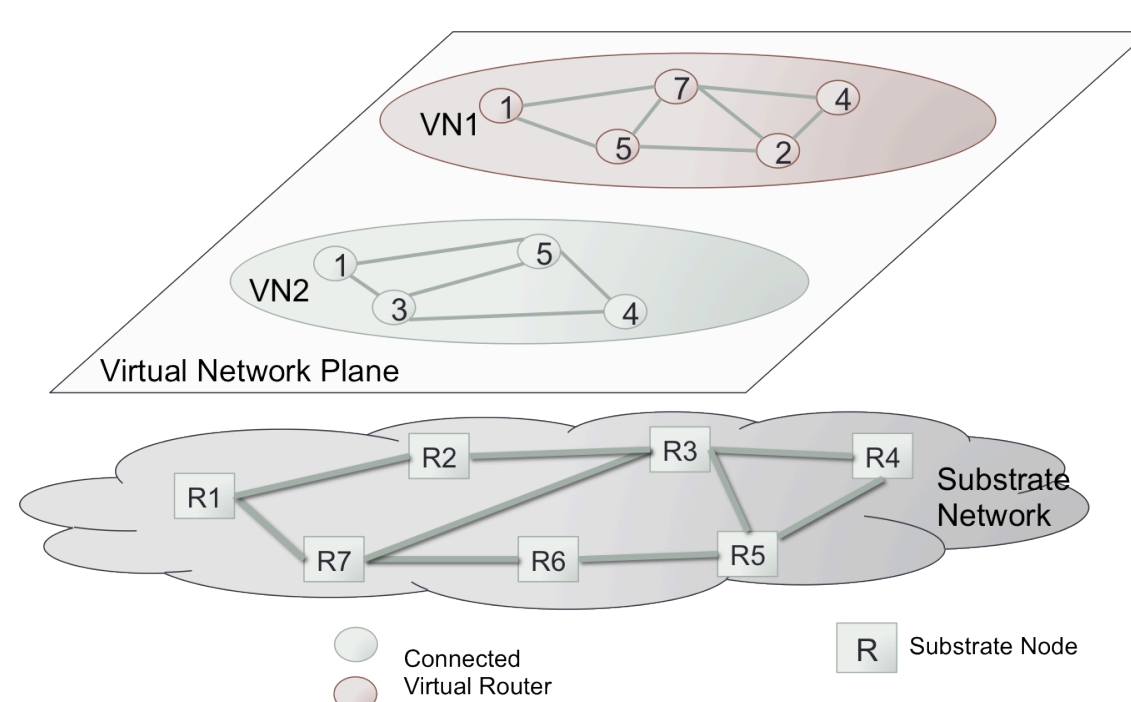
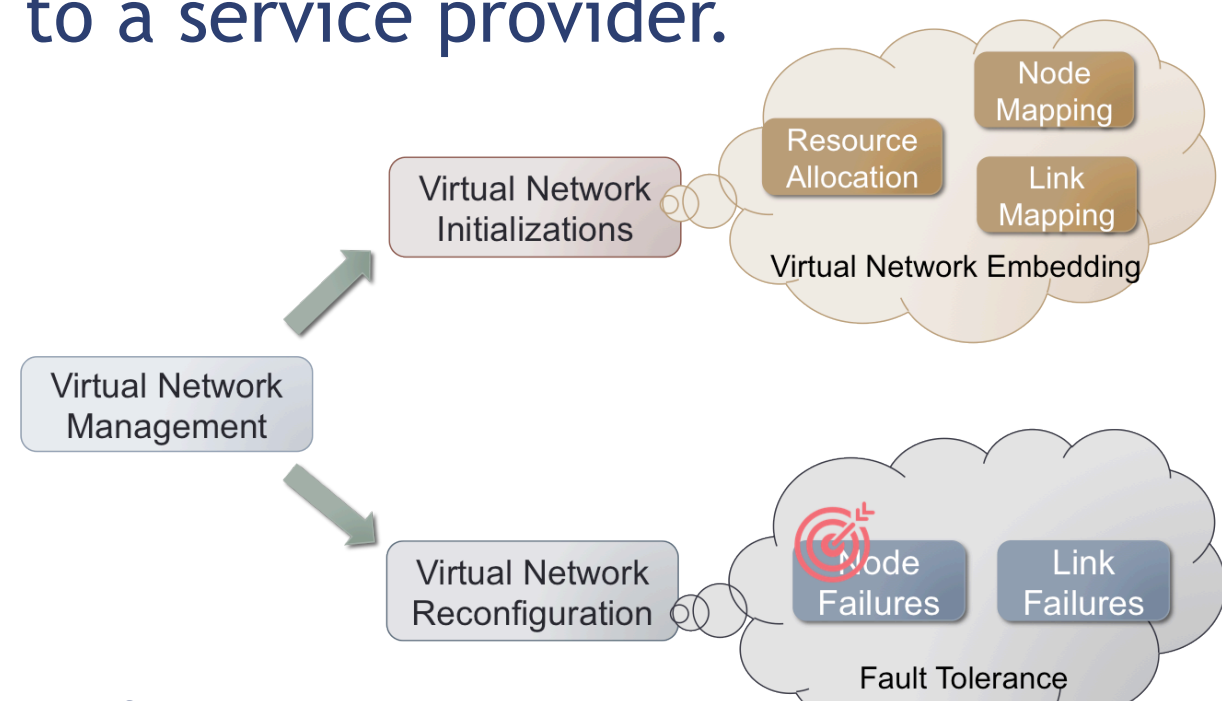


Fig. 1: Virtual Networks over the Substrate Network

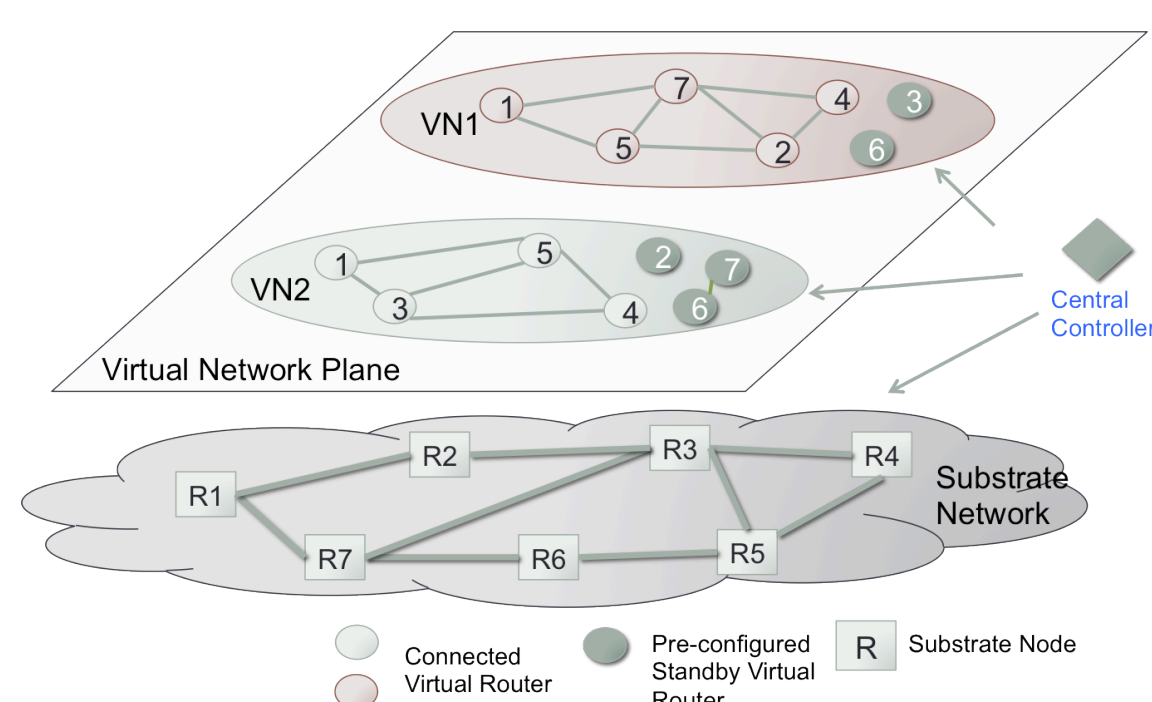


Fig. 2: Dynamic Reconfiguration Scheme for Virtual Network Environment

Our Goal: Optimal S-VR Selection Model

- Selects an optimal S-VR to restore single-VR failure for one or more VNs
 - The S-VR has optimal geographical location
 - Minimizing the maximum residual bandwidth utilization on the substrate nodes.
 - Subject to several main constraints
 - The VN topological structure does not change
 - The substrate node hosting the S-VR should have enough capacity

Assumptions

- Within a virtual network, any two virtual routers, including the standby virtual routers, are not hosted on the same substrate node.
- One substrate node failure at a time
- For each virtual network, at most one virtual router failure at a time

Type of Node Failures

- Independent Virtual Router Failure
 - Caused by software failures within the virtual nodes
 - Does not affect other virtual nodes in current or other virtual networks
- Dependent Virtual Router Failure
 - Caused by a failure occurring at a substrate node
 - Multiple virtual networks may get affected

Acknowledgement

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Optimization Model

Constraints

- Only a reserved S-VR can be used to restore a VR failure.

$$u_i^{j,f} \leq \delta_i^j (1 - \beta_i^j), \quad \forall i \in \mathcal{R}, \forall f \in \mathcal{F}^j, \forall j \in \mathcal{G}, i \neq f \quad (1)$$

- An S-VR can restore at most one VR failure in a VN

$$\sum_{f \in \mathcal{F}^j} \delta_i^j \cdot u_i^{j,f} \leq 1, \quad \forall i \in \mathcal{R}, \forall j \in \mathcal{G} \quad (3)$$

- Virtual Interfaces on an S-VR can be enabled only when this S-VR is selected.

$$\gamma_{i,m,k}^{j,f} \cdot v_{i,m,k}^{j,f} \leq u_i^{j,f}, \quad \forall f \in \mathcal{F}^j, \forall i, k \in \mathcal{R}, \forall j \in \mathcal{G}, \forall m \in \mathcal{M}_i^j, i \neq k \neq f \quad (5)$$

- Conservation of Virtual Link Bandwidth Requirement

$$\sum_{m \in \mathcal{M}_i^j} \sum_{k \in \mathcal{R}} \delta_i^j \cdot \gamma_{i,m,k}^{j,f} \cdot v_{i,m,k}^{j,f} \leq \sum_{p \in \mathcal{P}_i} \delta_i^j \cdot b_{i,p} \cdot u_i^{j,f}, \quad \forall j \in \mathcal{G}, \forall f \in \mathcal{F}^j, \forall i \in \mathcal{R}, i \neq k \neq f, \quad (7)$$

- Residual bandwidth utilization balancing on the substrate nodes.

$$\sum_{j \in \mathcal{G}} \sum_{f \in \mathcal{F}^j} \delta_i^j \cdot u_i^{j,f} \cdot \left(\sum_{k \in \mathcal{R}} \sum_{m \in \mathcal{M}_i^j} \gamma_{i,m,k}^{j,f} \cdot c_{i,m,k}^{j,f} \right) \leq b_{i,p} \cdot t, \quad \forall i \in \mathcal{R}, i \neq k \neq f \quad (8)$$

- Objective Function

$$\min \left\{ \lambda \cdot \sum_{j \in \mathcal{G}} \sum_{f \in \mathcal{F}^j} \sum_{i \in \mathcal{R}} \sum_{m \in \mathcal{M}_i^j} \sum_{k \in \mathcal{R}} \left(w_{\theta 1} \cdot \eta_{i,m,k}^{j,f} + w_{\theta 2} \cdot \sigma_{i,m,k}^{j,f} \right), \right. \\ \left. \delta_{i,j} \cdot \gamma_{i,m,k}^{j,f} \cdot v_{i,m,k}^{j,f} + \pi \cdot t, i \neq k \neq f \quad (11) \right.$$

Symbols	Description (all non-negative)
λ, π	Indicates the proportion for sub-objectives
w_{θ}	Primary weight parameters used in (11): $w_{\theta} = \langle w_{\theta 1}, w_{\theta 2} \rangle$
w_b	Secondary weights in (13): $w_b = \langle w_{b1}, w_{b2} \rangle$
w_a	Third level weights in (13): $w_a = \langle w_{a1}, w_{a2} \rangle$

- For any VR failure in a VN, only one S-VR can be selected

$$\sum_{i \in \mathcal{R}} \delta_i^j \cdot u_i^{j,f} = 1, \quad \forall f \in \mathcal{F}^j, \forall j \in \mathcal{G} \quad (2)$$

- The number of S-VRs selected from the same substrate node cannot exceed a predetermined limit h_i .

$$\sum_{f \in \mathcal{F}^j} \sum_{j \in \mathcal{G}} \delta_i^j \cdot u_i^{j,f} \leq h_i, \quad \forall i \in \mathcal{R} \quad (4)$$

- Conservation of VN topological Structures

$$\sum_{m \in \mathcal{M}_i^j} \sum_{k \in \mathcal{R}} e_{f,k}^j \cdot \delta_i^j \cdot \gamma_{i,m,k}^{j,f} \cdot v_{i,m,k}^{j,f} = \sum_{k \in \mathcal{R}} \delta_i^j \cdot u_i^{j,f} \cdot \beta_f^j \cdot e_{f,k}^j, \quad \forall j \in \mathcal{G}, \forall f \in \mathcal{F}^j, \forall i \in \mathcal{R}, i \neq k \neq f \quad (6)$$

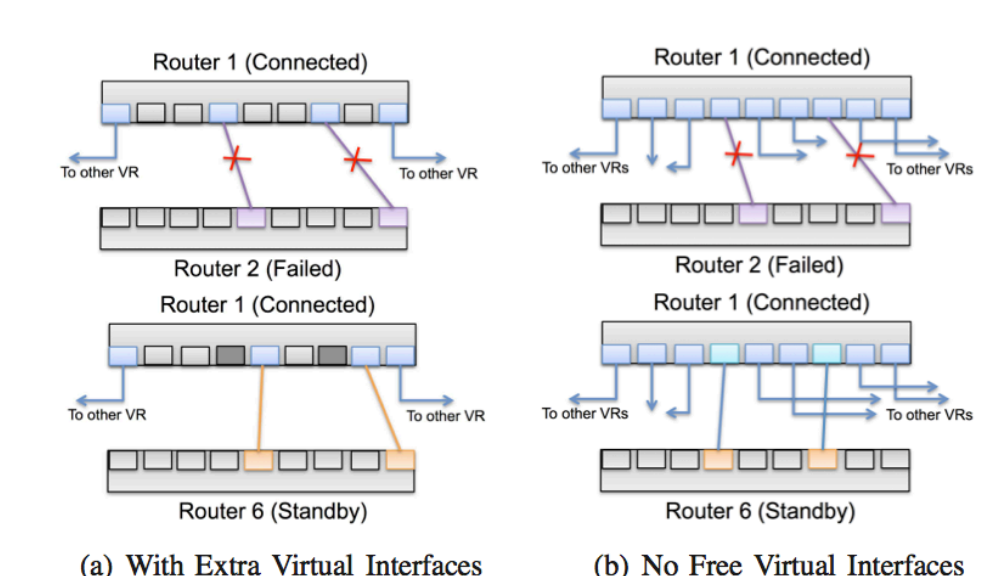
- Only select the S-VR when there is a path from its host to its neighbors' host on the substrate network.

$$\sum_{q \in \mathcal{Q}(i,k)} x_q^{(j,f,i,k)} = \sum_{m \in \mathcal{M}_i^j} \gamma_{i,m,k}^{j,f} \cdot c_{i,m,k}^{j,f} \cdot \delta_i^j \cdot u_i^{j,f}, \quad \forall j \in \mathcal{G}, \forall f \in \mathcal{F}^j, \forall i, k \in \mathcal{R}, i \neq k \neq f \quad (9)$$

$$\sum_{(j,f,i,k)} \sum_{q \in \mathcal{Q}(i,k)} x_q^{(j,f,i,k)} \leq b_i, \quad \forall i \in \mathcal{L}, i \neq k \neq f \quad (10)$$

- Cost Functions

- Type-I: Interface Operations



(a) With Extra Virtual Interfaces (b) No Free Virtual Interfaces

$$\eta_{i,m,k}^{j,f} = 2 \cdot (s^- + s^+) + \tau_i^{j,f} \cdot l, \quad \forall m \in \mathcal{M}_i^j, \forall k \in \mathcal{R}, i \neq k \neq f \quad (12)$$

- Type-II: Connectivity Cost

$$\sigma_{i,m,k}^{j,f} = w_{b1} \cdot (w_{a1} \cdot d_{i,f} + w_{a2} \cdot d_{i,k}) + w_{b2} \cdot rtt_{i,k}, \quad \forall j \in \mathcal{G}, \forall m \in \mathcal{M}_i^j, i \neq k \neq f \quad (13)$$

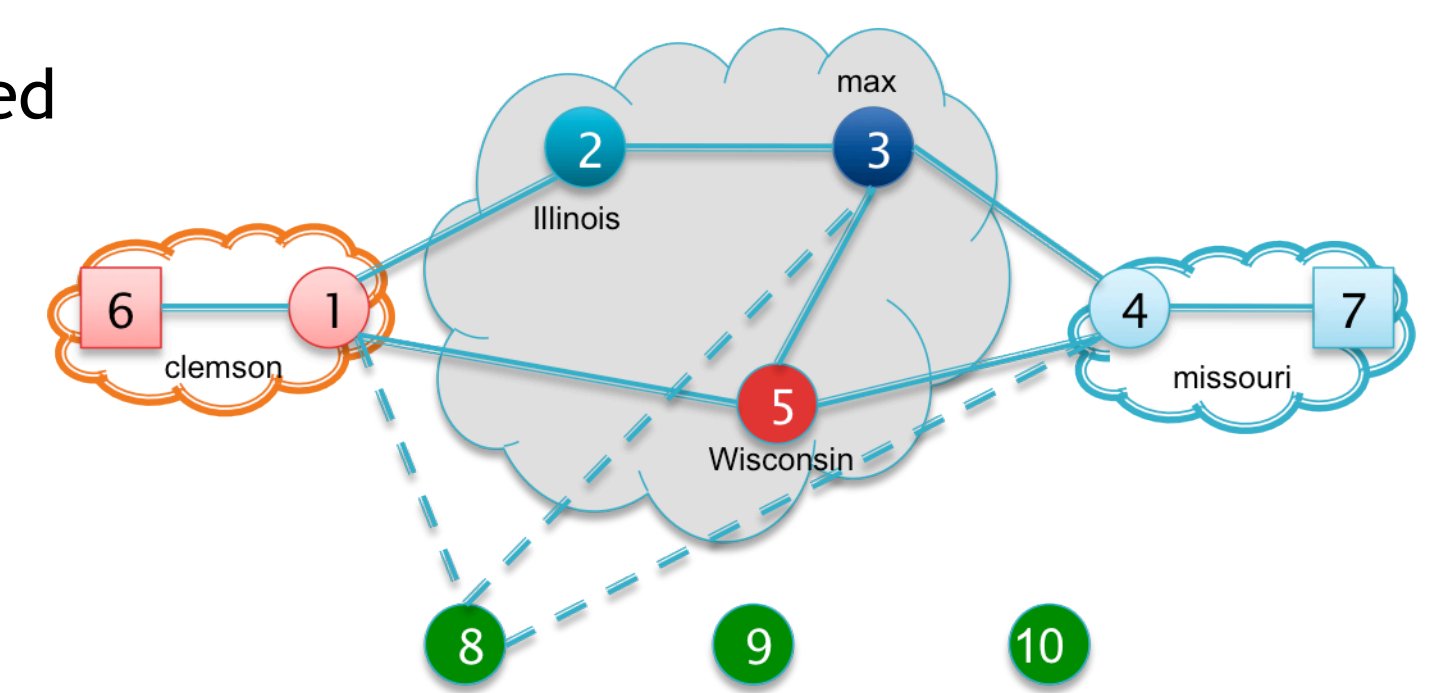
Preliminary Implementation on GENI

- Topology

- Seven Sites on GENI testbed
- One virtual network
- Three core virtual routers
- Three S-VR virtual routers
- One VR failure

- Selection Metrics

- Geographical location
- VM load on each VR's host.



SVR	Site	Latitude	Longitude	Dist (mile)	VM Load
8	Wisconsin	43.075	-89.41	0	6%
9	Stanford	37.430	-122.17	1760.078	14%
10	UtahDDC	40.751	-111.89	1164.032	4%

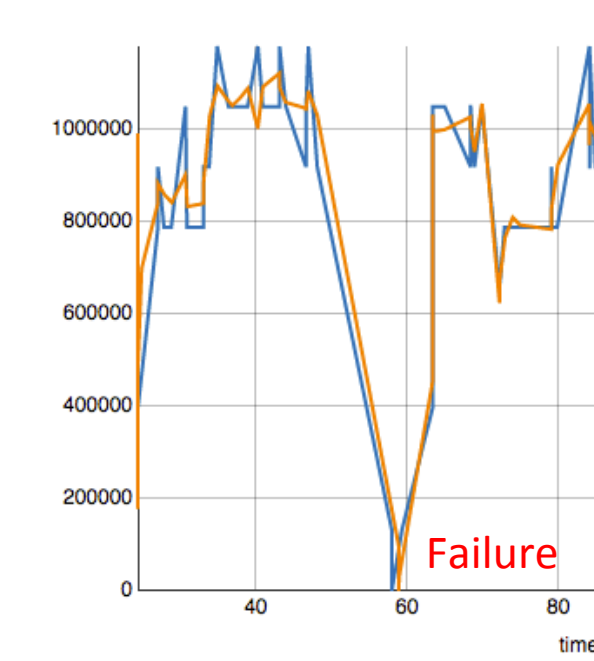


Figure: TCP throughput

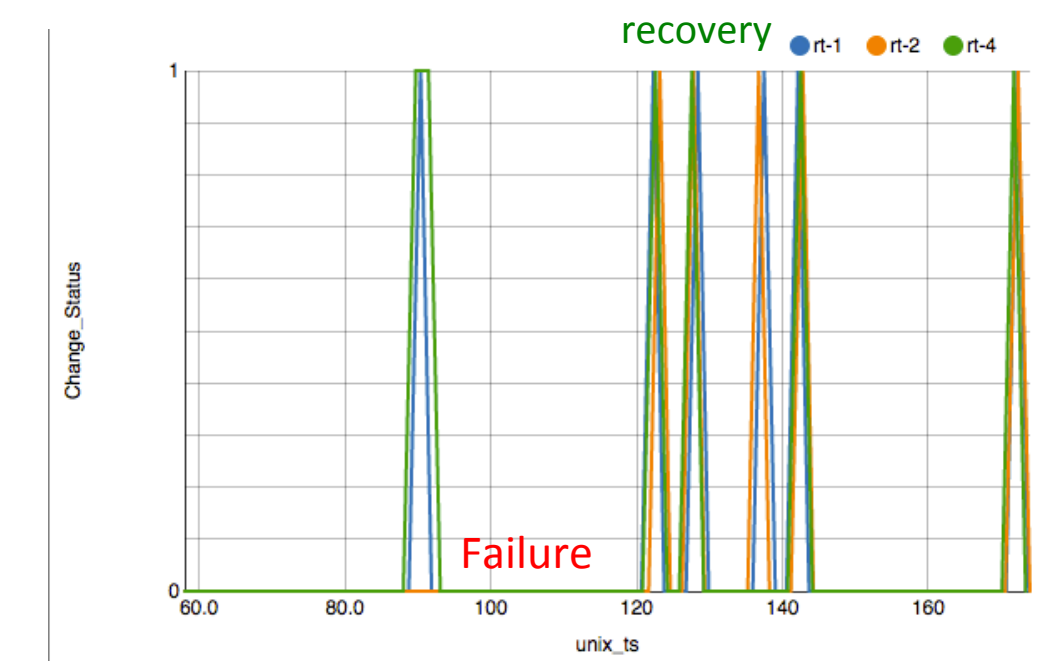


Figure: Routing Table Updates